

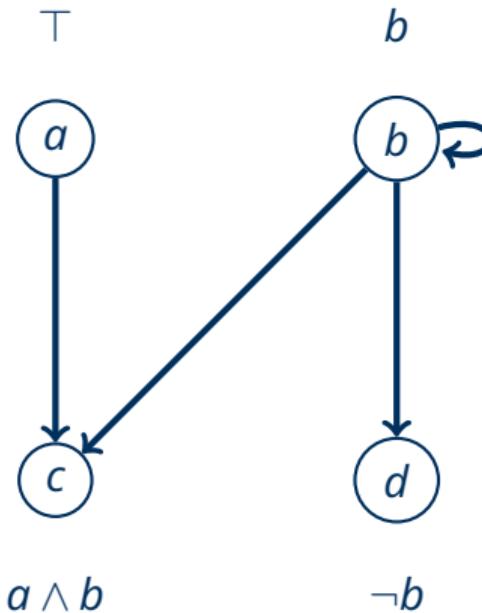
**Stefan Ellmauthaler, Sarah A. Gaggl, Dominik Rusovac, Johannes P. Wallner**

Faculty of Computer Science, International Center for Computational Logic, Knowledge-Based Systems Group

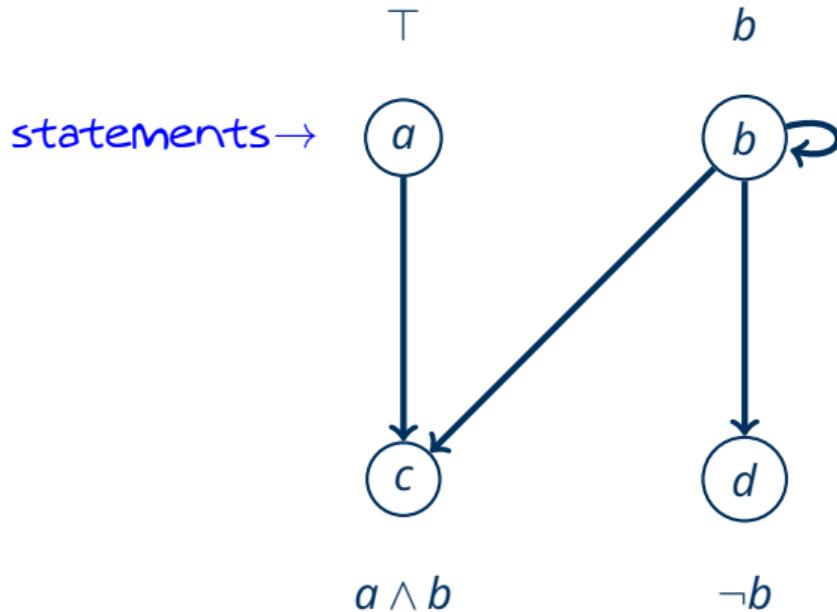
# ***Representing Abstract Dialectical Frameworks with Binary Decision Diagrams***

LPNMR 2022 // Genova Nervi, Italy, September 8, 2022

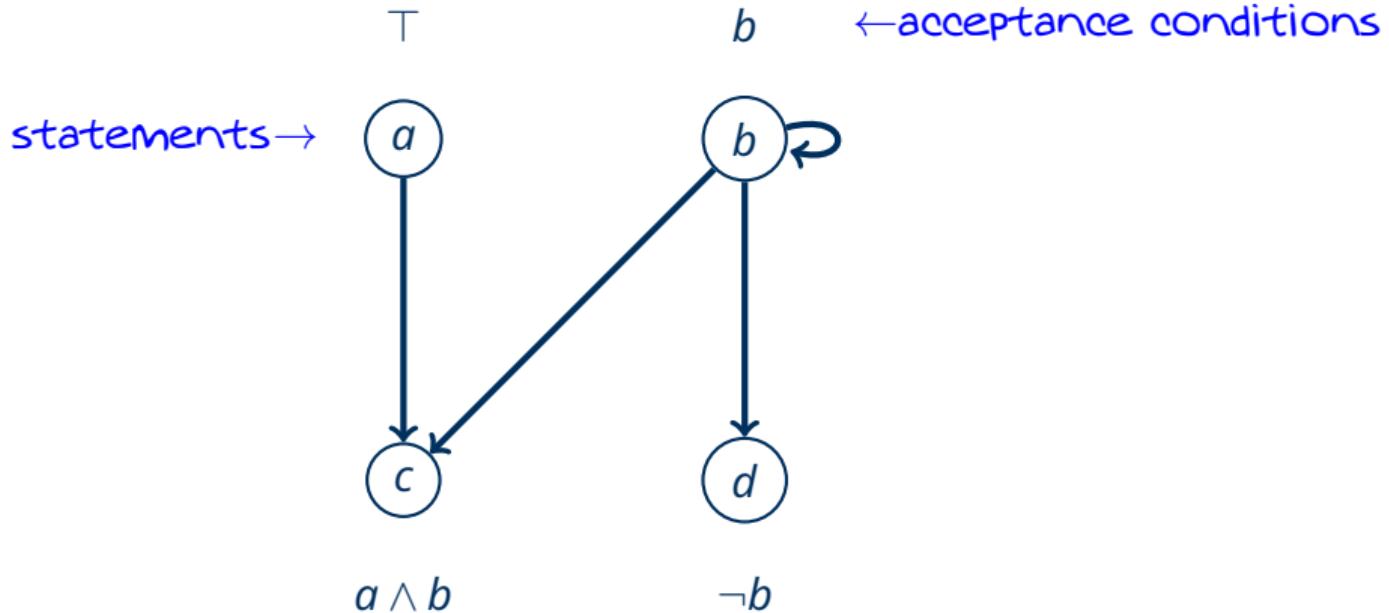
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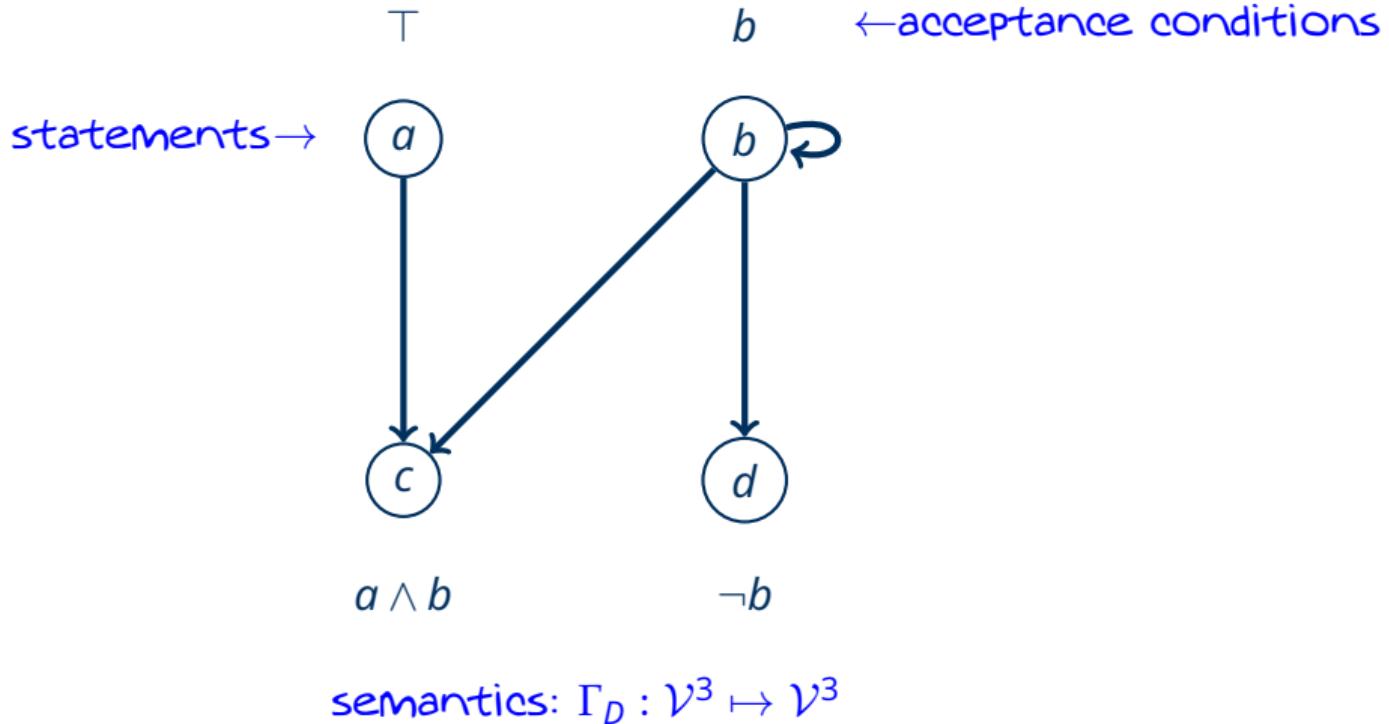
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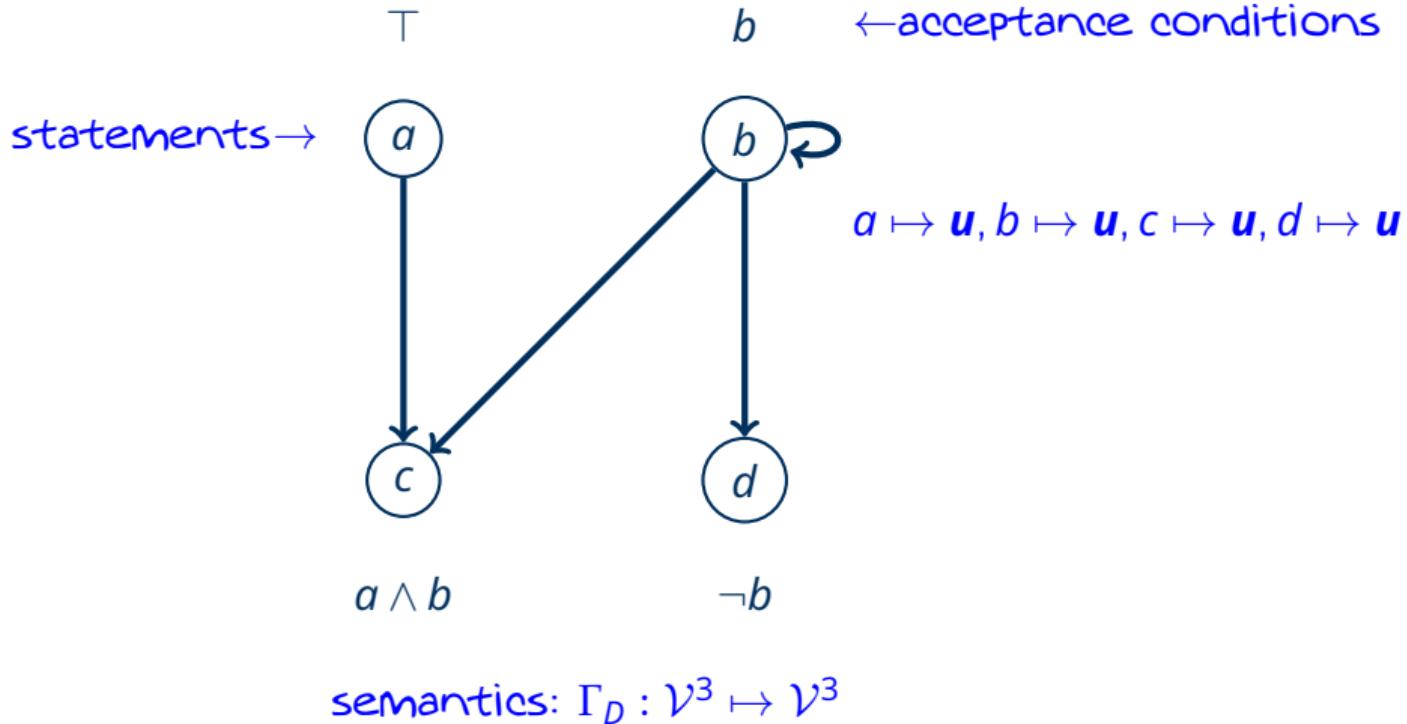
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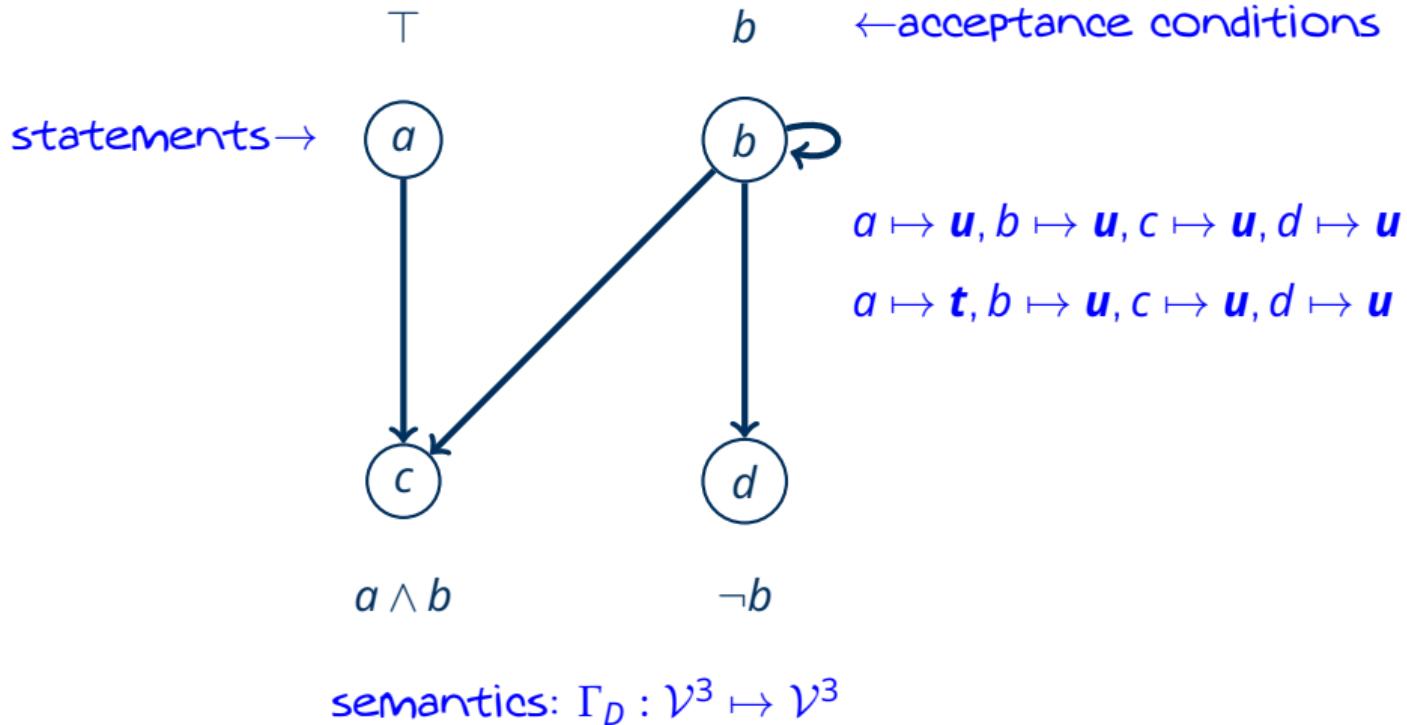
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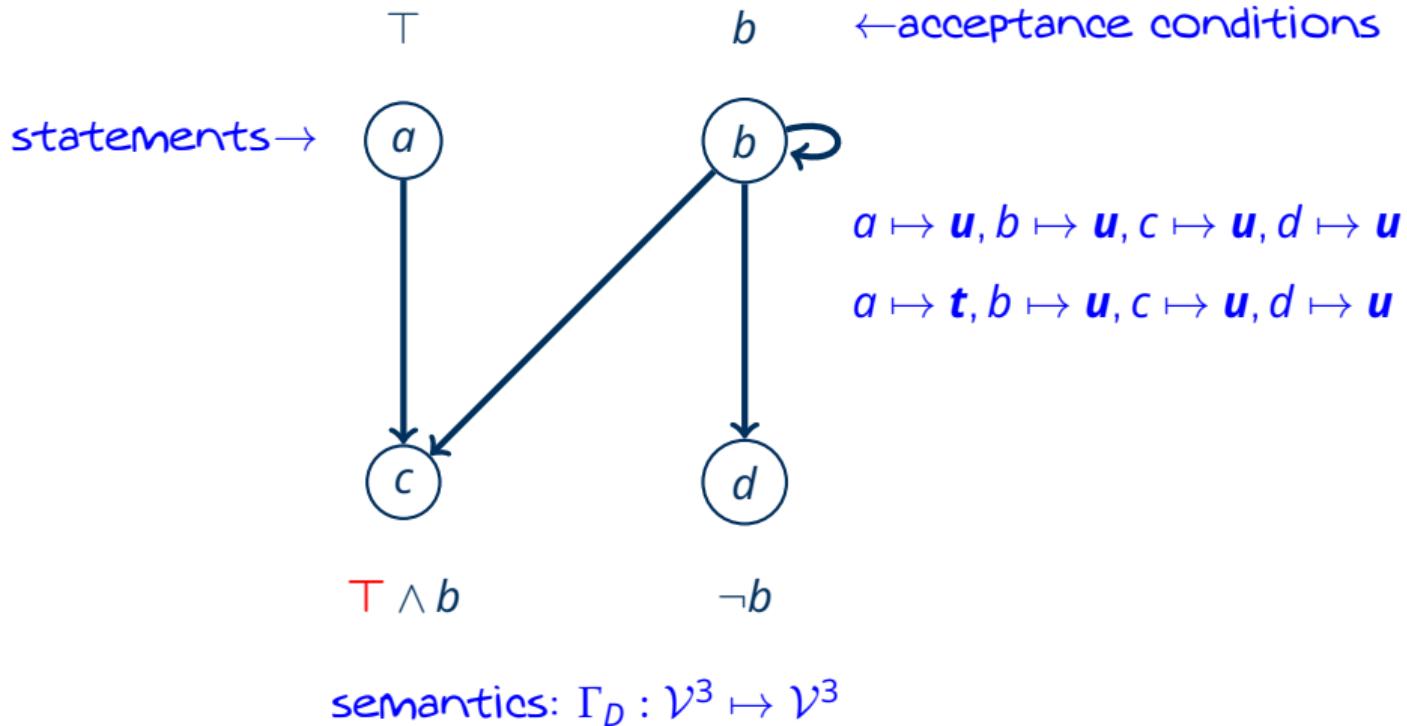
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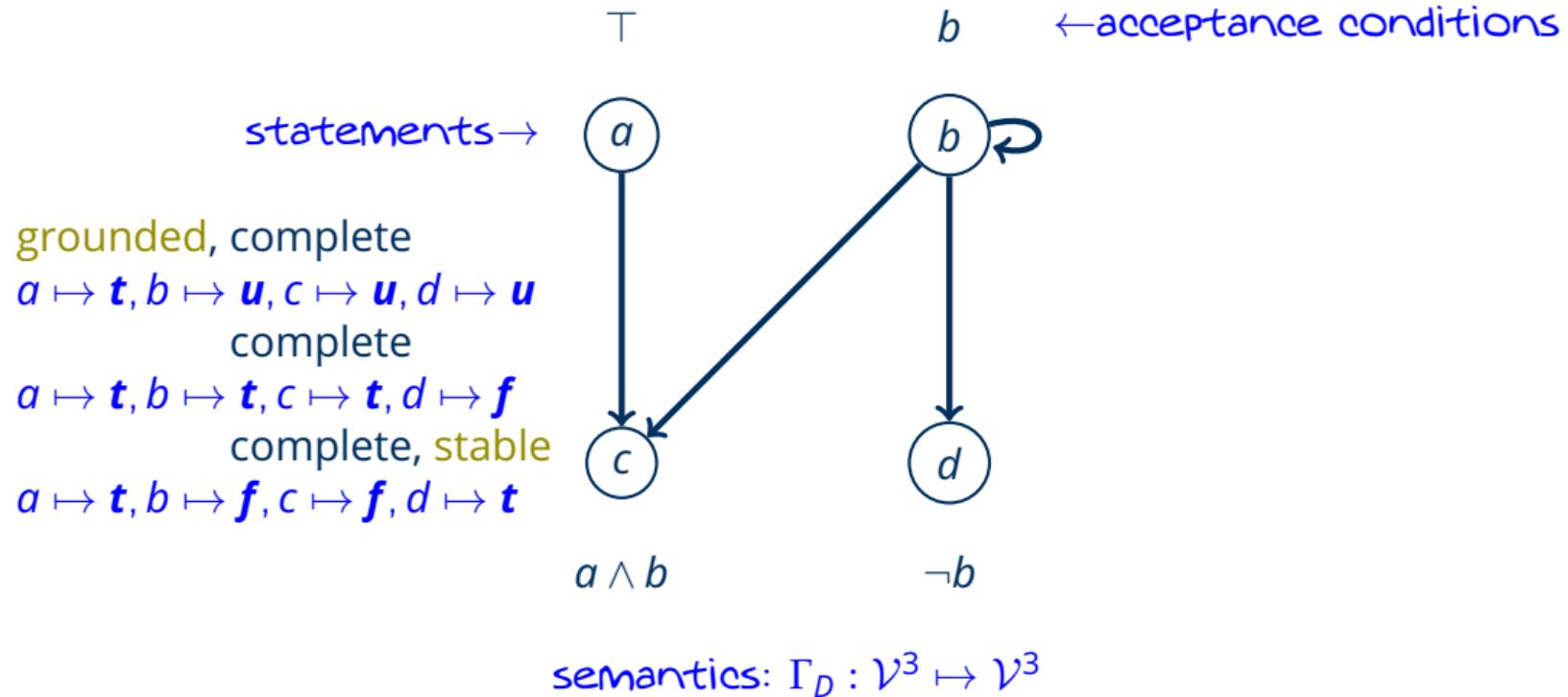
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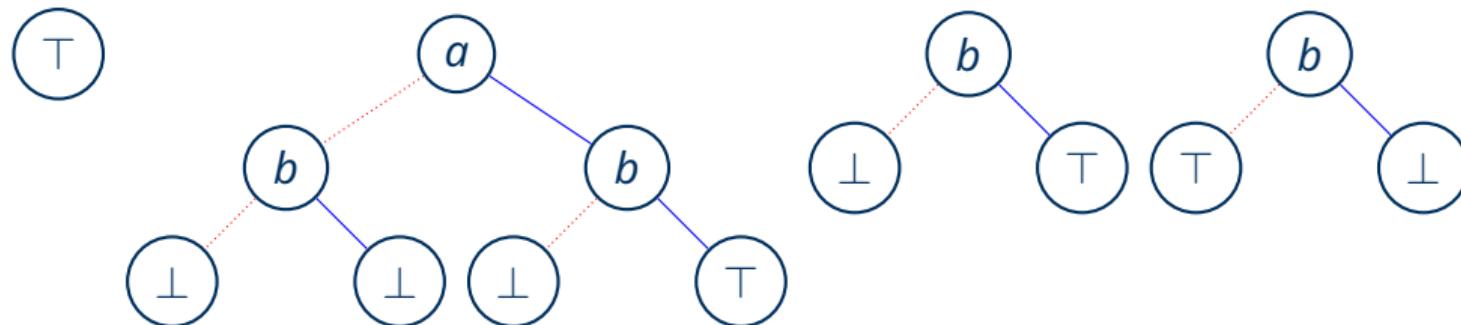


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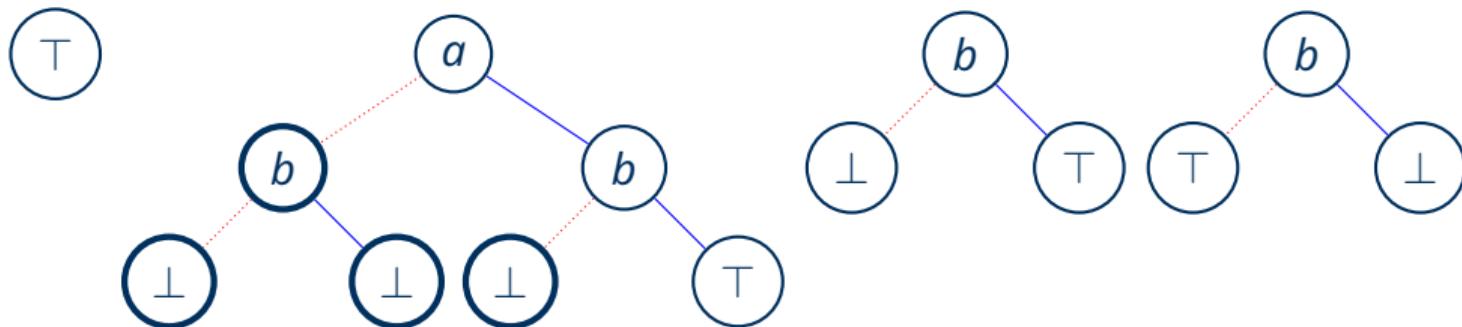
# ordered Binary Decision Tree

- Tree: inner nodes are variables and leafs are truth constants  $\top$  and  $\perp$
- Inner node has **lo** and **hi** child
- Every path from root to leaf needs to follow pre-defined strict ordering of variables



# reduced ordered Binary Decision Diagram [Bry86]

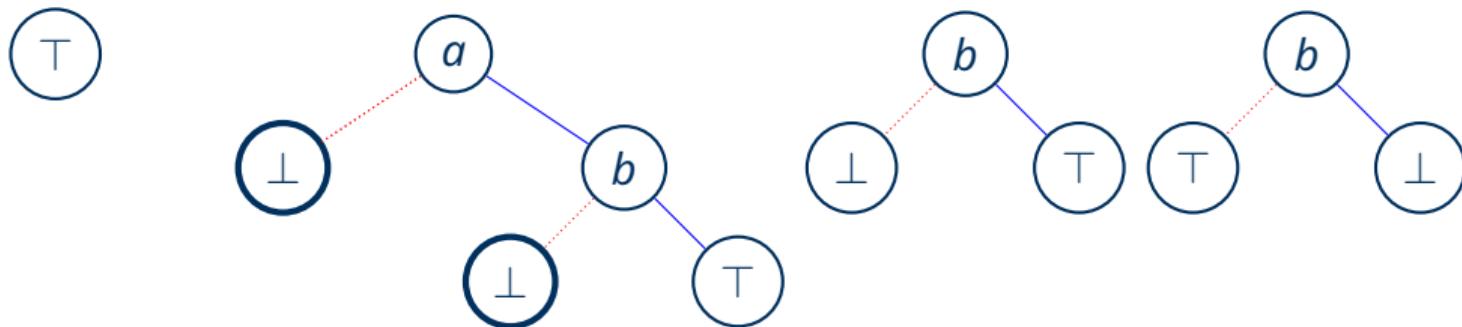
- $lo(n) = hi(n)$ , then replace  $n$  by  $hi(n)$
- if  $n = v$ , then replace  $v$  by  $n$  globally (violate tree-property)
- Given a variable order, this representation is unique under logical equivalence of formulae



# reduced ordered Binary Decision Diagram

## [Bry86]

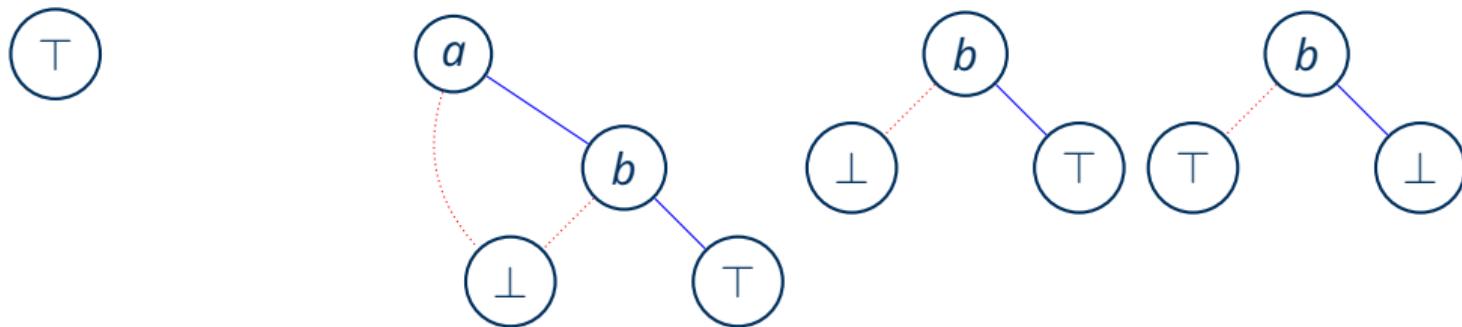
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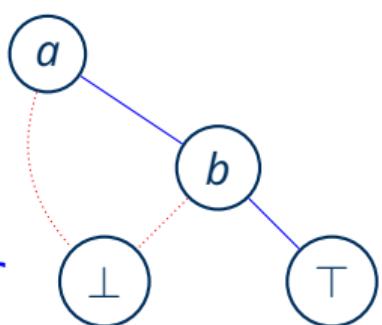


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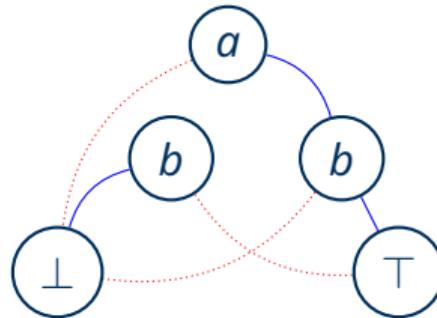
- Restriction linear



- Optimal variable order in NP
- Check for (un-)SAT and TAUT constant [DMO2]

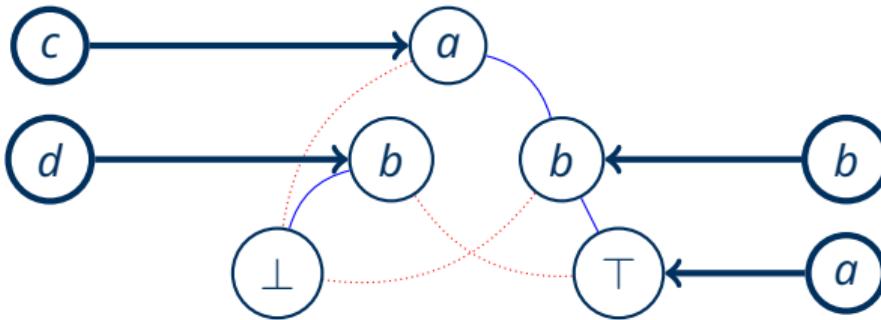
# New idea: roBDDs to represent ADFs

- To each statement, one BDD is related as the acceptance condition
- More compact representation due to "merging" of nodes



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# roBDDs to represent ADFs

## Theorem

*Given the BDD representation of an ADF  $D$ , the result of applying  $\Gamma_D$  to any three-valued interpretation  $\mathcal{I}$  can be computed in polynomial time.*

## Theorem

*Given an ADF  $D$  in BDD representation, there is a polynomial algorithm that computes the grounded interpretation of  $D$ .*

## Corollary

*Verifying whether a three-valued interpretation is a model or is stable in an ADF represented by BDDs is in P. Moreover credulous reasoning is in NP and sceptical reasoning in coNP.*

# ADF-BDD solver

- Written in Rust
- BDDs
  - own implementation
  - biodivine-bdd for faster instantiation
- Various BDD-modes (own, biodivine, hybrid)
- Grounded, complete, and stable semantics
- Github, Library, and Binary available
  - hub: <https://github.com/ellmau/adf-obdd>
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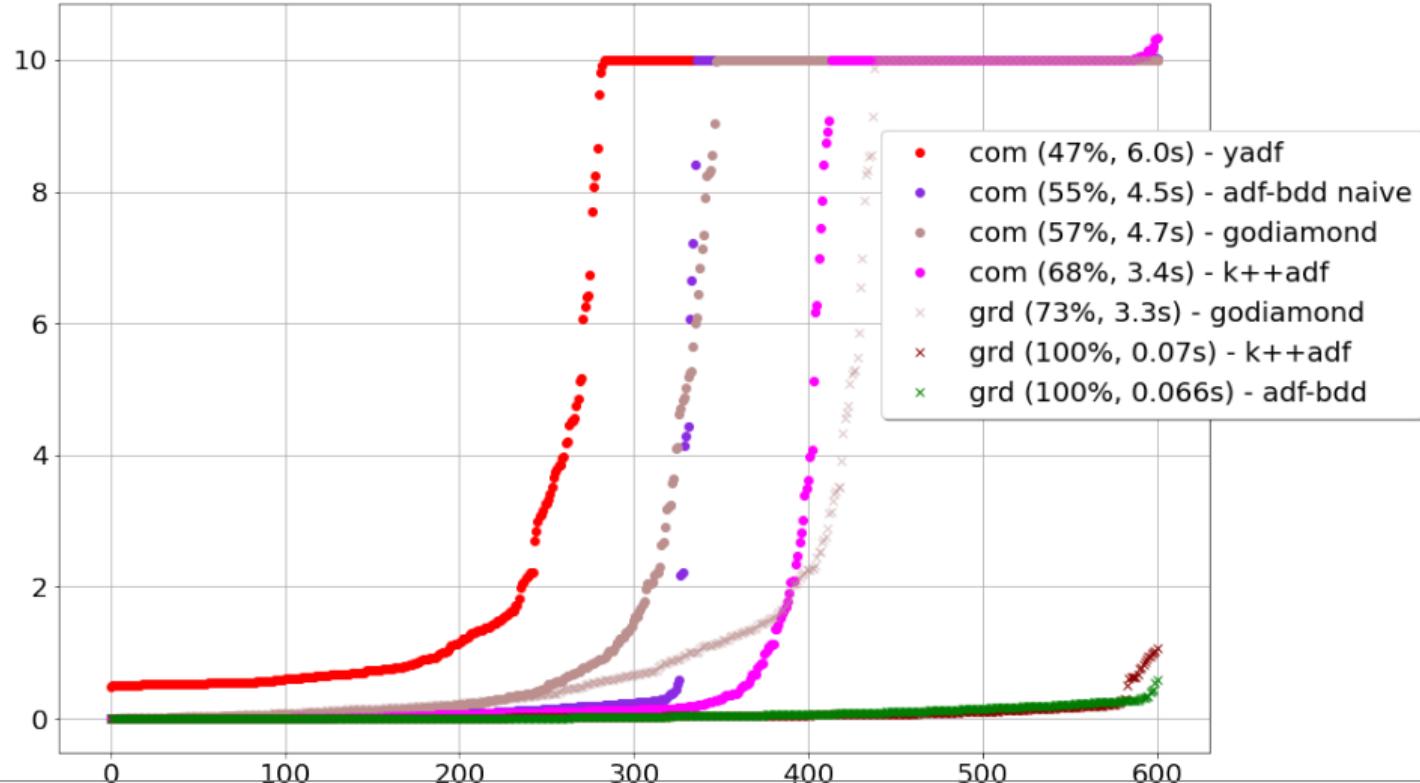
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## Evaluation

- goDIAMOND, k++adf, yadf
- 600 instances
- timeout 10 seconds
- **hyperfine evaluation**

# ADF-BDD Evaluation: grounded and complete



# Search Space Exploitation

## with Faceted Navigation

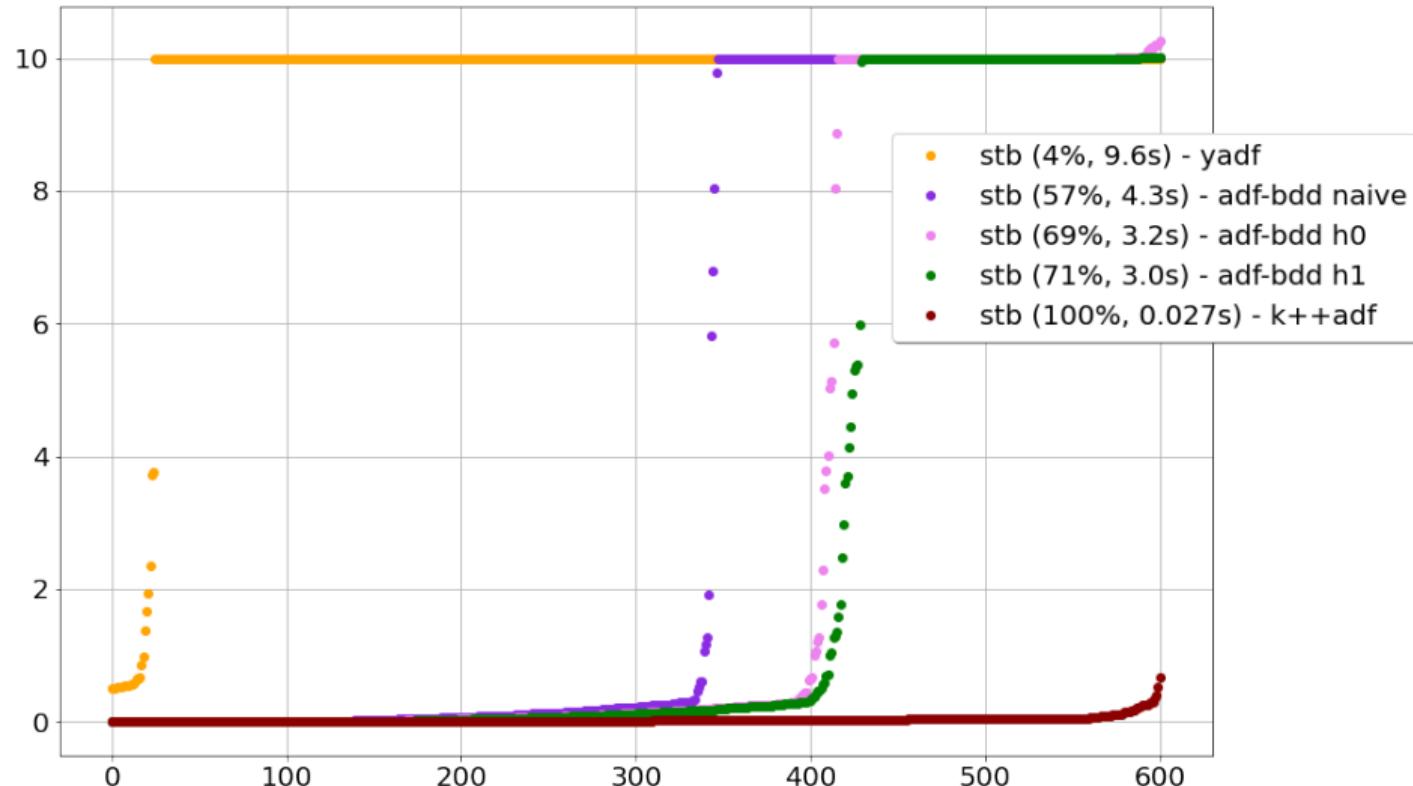
- Use Faceted Navigation measures to describe (Sub-)Search-Space
- Allows for an easy framework of properties for heuristics, like
  - Number of Models
  - Number of Facets
  - BDD Paths to  $\top$  resp.  $\perp$
  - Variable impact
  - ...
- Heuristics and Facet Navigation-based Algorithm for Stable Models
  - Recursive, one set of NoGood-like constraints per recursion path
  - Based on a heuristic, identify the optimal facet to activate
  - Propagates truth values, based on the facets and construct fixpoints
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# ADF-BDD Evaluation: stable



# Contributions

## ADF with BDD

- Use BDDs to represent ACs of ADFs
- Complexity analysis of this BDD-based ADFs
  - Drop of complexity - one level on polynomial hierarchy
  - Same as (easier) Dung AF and (less expressive) BADFs
- Unique representation of maximal information with respect to  $\Gamma_D$

## Facet Navigation for Search Space Exploitation

- Use Facet Navigation to navigate in the search space
- Represent Properties, Weights, and Heuristics in an uniform model

# Contributions II

## Software Tool

- Comparable to fastest solver for grounded semantics
- Comparable to 2nd fastest solver for complete semantics with a naive approach
- In between fastest and 2nd fastest for stable semantics with improved performance with faceted heuristics algorithm
- Easy usability through Rust-Ecosystem  
(e.g. `cargo install adf-bdd-bin` to try it out, lib-crate to use ADF-BDD in your own project,...)

# Future Work

- Implement full NoGood-reasoning
- Implement further Heuristics
- Improve the BDD methods
- Investigate optimal BDD-variable-orders
- Increase UX and Visualisation

# Thank you for your interest!

Visit ADF-BDD at

<https://ellmau.github.io/adf-obdd/>



# Abstract Dialectical Frameworks

## Definition (Abstract Dialectical Framework (ADF))

An ADF is a tuple  $\langle S, C \rangle$ , where

- $S$  is a fixed finite set of statements and
- $C := \{\varphi_s\}_{s \in S}$  is a set acceptance conditions for statements, which corresponds to propositional formulae whose variable signature is  $S$ .

# Abstract Dialectical Frameworks

## Definition ( $\Gamma_D$ )

Let  $D = \langle S, C \rangle$  be an ADF,  $\mathcal{I} : S \mapsto \{\mathbf{t}, \mathbf{f}, \mathbf{u}\}$  be a three-valued interpretation, and  $\Gamma_D(\mathcal{I}) : S \mapsto \{\mathbf{t}, \mathbf{f}, \mathbf{u}\}$  with

$$s = \begin{cases} \mathbf{t} & \text{if } \models \varphi_s(\mathcal{I}); \\ \mathbf{f} & \text{if } \varphi_s(\mathcal{I}) \models \perp; \\ \mathbf{u} & \text{otherwise.} \end{cases}$$

## Definition (Semantics)

$\mathcal{I}$  is

- **complete** if  $\mathcal{I} = \Gamma_D(\mathcal{I})$
- **grounded** if  $\mathcal{I} = \text{lfp}(\Gamma_D)$

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- **stable** if  $\mathcal{I}$  is two-valued, complete, and for the grounded interpretation  $\mathcal{W}$  of  $D^{\mathcal{I}}$  it holds that  $\mathcal{I}(s) = \mathbf{t}$  implies  $\mathcal{W}(s) = \mathbf{t}$

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## Definition (Reduction)

Let  $D = \langle S, C \rangle$  be an ADF,  $\mathcal{I} : S \mapsto \{\mathbf{t}, \mathbf{f}\}$  be a two-valued interpretation.  
 $D^{\mathcal{I}} = \langle S^{\mathcal{I}}, C^{\mathcal{I}} \rangle$ , where

- $S^{\mathcal{I}} = \{s \in S \mid \mathcal{I}(s) = \mathbf{t}\}$
- $C^{\mathcal{I}} = \{\varphi_s[s'/\perp : \mathcal{I}(s') = \mathbf{f}] \}$

# References I

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